# Improved Abramov-Petkovšek's Reduction for Hypergeometric Terms 

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## Hypergeometric Sequences

A sequence $T: \mathbb{N} \rightarrow \mathbb{C}$ is said to be hypergeometric if $\exists R \in \mathbb{C}(n)$ s.t. $T(n+1)=R(n) T(n)$ for all $n \gg 0$.


## Notation.

- For a sequence $A(n), \Delta(A)=A(n+1)-A(n)$.
- For two hypergeometric sequences $T(n), S(n)$, write

$$
T \equiv S \text { if } T-S=\Delta(H) \text { for some hypergeometric sequence } H(n) \text {. }
$$

Definition 1. A hypergeometric sequence $T(n)$ is summable if $T \equiv 0$.
Summability. Given a hypergeometric sequence $T(n)$, decide whether $T \equiv 0$.
Gosper algorithm [4] solves the summability problem.

## Abramov-Petkovšek's Reduction

Additive Decomposition. Given a hypergeometric sequence $T(n)$, compute two hypergeometric sequences $T_{1}(n)$ and $T_{2}(n)$ s.t.

$$
\begin{equation*}
T=\Delta\left(T_{1}\right)+T_{2} \quad \text { and } \quad T \equiv 0 \Longleftrightarrow T_{2}=0 \tag{1}
\end{equation*}
$$

Abramov-Petkovšek's reduction (AP reduction) [1, 2] computes $T_{1}(n)$ and $T_{2}(n)$ in (1). The AP reduction solves the summability problem as well.

## Definition 2.

$-p \in \mathbb{C}[n]$ is shift-free if $\operatorname{gcd}(p(n), p(n+i))=1$ for all $i \in \mathbb{Z} \backslash\{0\}$
$-r \in \mathbb{C}(n)$ with $r=u / v$ is shift-reduced if $\operatorname{gcd}(u(n), v(n+i))=1$ for all $i \in \mathbb{Z}$.

- Let $r=u / v$ be a shift-reduced rational function in $\mathbb{C}(n)$. A polynomial $f \in \mathbb{C}[n]$ is strongly prime with $r$ if either $f \in \mathbb{C}$, or, for every irreducible factor $p$ of $f$,

$$
p \nmid u v, p(n+i) \nmid u \text { and } p(n-i) \nmid v \text { for all } i \in \mathbb{Z}^{+} .
$$

AP reduction. Given a hypergeometric sequence $T(n)$, compute two hypergeometric sequences $T_{1}(n)$ and $T_{2}(n)$ such that the two conditions in (1) hold.

1. Compute three polynomials $a, b, w \in \mathbb{C}[n]$ such that

$$
\begin{equation*}
T \equiv\left(\frac{a}{b}+\frac{w}{v}\right) H \tag{2}
\end{equation*}
$$

where $H(n+1) / H(n)=u / v$ is shift-reduced, $\operatorname{deg}(a)<\operatorname{deg}(b), \operatorname{gcd}(a, b)=1$ and $b$ is shift-free and strongly prime with $u / v$. Moreover, the degree of the numerator $w$ is bounded
2. Consider the equation
$u(n) y(n+1)-v(n) y(n)=w(n)$.
(3)

Summable case
$b \in \mathbb{C}^{*}$ and $w=0$ or
(3) has a polynomial solution

$$
\begin{gathered}
\Downarrow \\
T_{2}=0
\end{gathered}
$$

Non-summable case
$b \notin \mathbb{C}^{*}$ or
(3) has no polynomial solution
$T_{2}=(a / b+w / v) H$
The sequence $T_{1}(n)$ can be constructed as the product of a rational function $r(n)$ and the sequence $H(n)$ incrementally.

## Improved AP Reduction

Idea. Not only bound the degree of the numerator $w$ in (2) as in [1, 2], but also reduce the number of its terms as in [3].
Definition 3. Let $K=u / v$ be shift-reduced. Define

$$
\begin{aligned}
\phi_{K}: \mathbb{C}[n] & \longrightarrow \mathbb{C}[n] \\
f(n) & \mapsto u(n) f(n+1)-v(n) f(n) .
\end{aligned}
$$

Let $\mathcal{N}_{K}=\operatorname{span}_{\mathbb{C}}\left\{n^{\ell} \mid \ell \in \mathbb{N}\right.$ and $\ell \neq \operatorname{deg}(g)$ for all $\left.g \in \operatorname{im}\left(\phi_{K}\right)\right\}$

Key Lemma. The $\mathbb{C}$-linear map $\phi_{K}$ is injective and $\mathbb{C}[n]=\operatorname{im}\left(\phi_{K}\right) \oplus \mathcal{N}_{K}$
Improved AP reduction. Given a hypergeometric sequence $T(n)$, compute two hypergeometric sequences $T_{1}(n)$ and $T_{2}(n)$ s.t. the two conditions in (1) hold.

## 1. Compute (2) as in step 1 of AP reduction.

2. Compute the projection $p$ of $w$ in $\mathcal{N}_{K}$. Set

$$
T_{2}:=\left(\frac{a}{b}+\frac{p}{v}\right) H
$$

The sequence $T_{1}(n)$ can be constructed as the product of a rational function $r(n)$ and the sequence $H(n)$ incrementally.

The improved AP reduction avoids computing a polynomial solution of any auxiliary $\mathrm{O} \Delta \mathrm{E}$.

## Experiments

We compare

- G: the Maple function Gosper in SumTools[Hypergeometric]
- S: a procedure that solves the summability problem based on improved AP, in which $T_{1}$ is not normalized if $T$ is not summable.
- AP: the Maple function SumDecomposition in SumTools[Hypergeometric]
- IAP: the reduction algorithm of improved AP-reduction, in which $T_{1}$ is always normalized

Test suite:

$$
T(n):=\frac{f(n)}{g(n) \cdot g(n+\lambda) \cdot g(n+\mu) \cdot h(n) \cdot h(n+\lambda) \cdot h(n+\mu)} \cdot \prod_{k=n_{0}}^{n} \frac{u(k)}{v(k)},
$$

where $f, g, h \in \mathbb{Z}[n]$ of respective degrees 20,10 and $10, u(n), v(n)$ are the product of two linear polynomials in $\mathbb{Z}[n]$, and $\lambda, \mu \in \mathbb{N}$ with $\lambda \leq \mu$.

| Input: $T(n)$ |  |  |  | Input: $T(n+1)-T(n)$ |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $(\lambda, \mu)$ | G S | AP | IAP | $(\lambda, \mu)$ | G S | AP | IAP |
| [0, 0] | 0.080 .12 | 0.19 | 0.12 | [0, 0] | 1.221 .46 | 2.83 | 1.44 |
| [5, 5] | 0.420 .52 | 4.80 | 0.64 | [5, 5] | 1.981 .75 | 9.06 | 1.76 |
| [10, 10] | 0.741 .00 | 17.06 | 1.42 | [10, 10] | 2.551 .87 | 19.21 | 1.89 |
| [10, 20] | $3.05 \quad 2.08$ | 66.50 | 4.30 | [10, 20] | 6.112 .55 | 49.43 | 2.55 |
| [10, 30] | 9.183 .53 | 237.50 | 10.54 | [10, 30] | 16.272 .66 | 111.77 | 2.70 |
| [10, 40] | 20.385 .20 | 482.34 | 24.02 | [10, 40] | 31.562 .88 | 214.57 | 2.89 |

Timings (in sec.) measured on a Mac computer, 4GB RAM, 3.06 GHz Core 2 Duo processor.


Experiments illustrate that the improved AP reduction is more efficient than both Gosper algorithm and AP-reduction.

## A Potential Application

Can one compute the minimal telescoper for a bivariate hypergeometric term by the improved AP reduction, following the idea in [3]?
Advantage. Such an algorithm would separate the computation for telescopers from that for certificates so as to improve efficiency
Difficulty. The least common multiple of shift-free polynomials is not necessarily shift-free.

## References

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