FOUNTAINS, HISTOGRAMS, AND q–IDENTITIES

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ABSTRACT. We solve the recursion $S_n = S_{n-1} - q^n S_{n-p}$, both, explicitly, and in the limit for $n \to \infty$, proving in this way a formula due to Merlini and Sprugnoli. It is also discussed how computer algebra could be applied.

1. Generalized Schur polynomials

In the paper [5], a short and direct proof of the following fact would be required [4]: For fixed $p \geq 1$, define polynomials as follows:

$$E_n = E_{n-1} - q^n E_{n-p}, \quad n \geq p, \qquad E_0 = \cdots = E_{p-1} = 1,$$

$$D_n = D_{n-1} - q^n D_{n-p}, \quad n \geq p, \qquad D_i = 1 - \sum_{j=1}^{i} q^j, \ i = 0, \ldots, p-1.$$

Then

$$E_\infty := \lim_{n\to\infty} E_n = \sum_{k\geq 0} \frac{(-1)^k q^{p\binom{k+1}{2}}}{(1-q)\ldots(1-q^k)},$$

$$D_\infty := \lim_{n\to\infty} D_n = \sum_{k\geq 0} \frac{(-1)^k q^{k+p\binom{k}{2}}}{(1-q)\ldots(1-q^k)}.$$

We will not only achieve that but actually derive *explicit* expressions for these polynomials!

Let us study the generic recursion

$$S_n = S_{n-1} + tq^{n-p} S_{n-p},$$

with unspecified initial values $S_0, \ldots, S_{p-1}$. For $p = 2$, these polynomials were studied by Andrews (and others) in the context of *Schur polynomials*, see [2].

We will use standard notation from q–calculus, see [1]:

$$(x)_n = (1-x)(1-xq)\ldots(1-xq^{n-1}), \qquad \begin{bmatrix} n \\ k \end{bmatrix} = \frac{(q)_n}{(q)_k (q)_{n-k}}.$$

It will be convenient to define $\begin{bmatrix} n \\ k \end{bmatrix} = 0$ for $n < 0$ or $k > n$.

Now we will proceed as in [1] and consider noncommutative variables x, η, such that $x\eta = q\eta x$; all other variables commute.

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Lemma 1.

$$\big(x+x^p\eta\big)^n=\sum_{k=0}^{n}\begin{bmatrix}n\\k\end{bmatrix}q^{p\binom{n}{2}-pnk+p\binom{k+1}{2}}x^{k+p(n-k)}\eta^{n-k}.$$

Proof. We write

$$\big(x+x^p\eta\big)^n=\sum_{k=0}^{n}a_{n,k}x^{k+p(n-k)}\eta^{n-k},$$

and $\big(x+x^p\eta\big)^{n+1}=\big(x+x^p\eta\big)^n\big(x+x^p\eta\big)$ resp. as $\big(x+x^p\eta\big)^{n+1}=\big(x+x^p\eta\big)\big(x+x^p\eta\big)^n$, compare coefficients, and get the recursions

$$\begin{aligned}a_{n+1,k}&=a_{n,k-1}+a_{n,k}q^{k+p(n-k)},\\ a_{n+1,k}&=a_{n,k-1}q^{n+1-k}+a_{n,k}q^{p(n-k)}.\end{aligned}$$

From this we derive, taking differences,

$$a_{n,k}=\frac{1-q^{n+1-k}}{1-q^k}q^{-p(n-k)}a_{n,k-1}.$$

The result follows from iteration by noting that $a_{n,0}=q^{p\binom{n}{2}}$. □

Of course we also have

$$\big(x+tx^p\eta\big)^n=\sum_{k=0}^{n}\begin{bmatrix}n\\k\end{bmatrix}q^{p\binom{n}{2}-pnk+p\binom{k+1}{2}}x^{k+p(n-k)}t^{n-k}\eta^{n-k}.$$

Now we derive the generating function for

$$F(x)=\sum_{n\ge0}S_nx^n;$$

the following procedure is inspired by [2]. Note that we can alternatively view η as an operator, defined by $\eta f(x)=f(qx)$. Cigler worked also much with this technique [3]. We find

$$\sum_{n\ge p}S_nx^n=\sum_{n\ge p}S_{n-1}x^n+\sum_{n\ge p}tq^{n-p}S_{n-p}x^n=x\sum_{n\ge p-1}S_nx^n+tx^p\sum_{n\ge0}\eta S_nx^n$$

or

$$F(x)-\sum_{n<p}S_nx^n=xF(x)-x\sum_{n<p-1}S_nx^n+tx^p\eta F(x),$$

and

$$F(x)=\frac{1}{1-x-tx^p\eta}\bigg(\sum_{i<p}S_ix^i-\sum_{i<p-1}S_ix^{i+1}\bigg).$$

Now we can apply our lemma and write

$$\begin{aligned}F(x)&=\sum_{n\ge0}\big(x+tx^p\eta\big)^n\bigg(\sum_{i<p}S_ix^i-\sum_{i<p-1}S_ix^{i+1}\bigg)\\
&=\sum_{n\ge0}\sum_{k=0}^{n}\begin{bmatrix}n\\k\end{bmatrix}q^{p\binom{n}{2}-pnk+p\binom{k+1}{2}}x^{k+p(n-k)}t^{n-k}\eta^{n-k}\bigg(\sum_{i<p}S_ix^i-\sum_{i<p-1}S_ix^{i+1}\bigg)\\
&=\sum_{n\ge0}\sum_{k=0}^{n}\begin{bmatrix}n\\k\end{bmatrix}q^{p\binom{n}{2}-pnk+p\binom{k+1}{2}}x^{k+p(n-k)}t^{n-k}\bigg(\sum_{i<p}S_iq^{i(n-k)}x^i-\sum_{i<p-1}S_iq^{(i+1)(n-k)}x^{i+1}\bigg)\end{aligned}$$

$$= \sum_{n\ge 0}\sum_{k=0}^{n} \begin{bmatrix} n\\ k\end{bmatrix} q^{p\binom{k}{2}} x^{n-k+pk} t^k \Bigg(\sum_{i<p} S_i q^{ik} x^i - \sum_{i<p-1} S_i q^{(i+1)k} x^{i+1} \Bigg)$$
$$= \sum_{k,n\ge 0} \begin{bmatrix} n+k\\ k\end{bmatrix} q^{p\binom{k}{2}} x^{n+pk} t^k \Bigg(\sum_{i<p} S_i q^{ik} x^i - \sum_{i<p-1} S_i q^{(i+1)k} x^{i+1} \Bigg)$$
$$= \sum_{k\ge 0} q^{p\binom{k}{2}} x^{pk} t^k \frac{1}{(x)_{k+1}} \Bigg(\sum_{i<p} S_i q^{ik} x^i - \sum_{i<p-1} S_i q^{(i+1)k} x^{i+1} \Bigg).$$

From this we find an explicit formula for S_n (the quantity S_{-1} has to be interpreted as 0):

$$S_n = \sum_{0\le i<p} (S_i - S_{i-1}) \sum_{k\ge 0} \begin{bmatrix} n-(p-1)k-i\\ k\end{bmatrix} q^{p\binom{k}{2}+ik} t^k.$$

Now we specialize this to our instance. Here, $t=-q^p$, and thus

$$S_n = \sum_{0\le i<p} (S_i - S_{i-1}) \sum_{k\ge 0} \begin{bmatrix} n-(p-1)k-i\\ k\end{bmatrix} q^{p\binom{k+1}{2}+ik} (-1)^k.$$

Therefore

$$E_n = \sum_{k\ge 0} \begin{bmatrix} n-(p-1)k\\ k\end{bmatrix} q^{p\binom{k+1}{2}} (-1)^k.$$

From this, the limit of E_n is immediate. For D_n we eventually get the following form

$$D_n = \sum_{k\ge 0} \begin{bmatrix} n-(p-1)(k-1)\\ k\end{bmatrix} q^{k+p\binom{k}{2}} (-1)^k,$$

from which the formula for D_∞ is immediate. To prove it, we need a simple lemma whose proof is just a routine calculation.

Lemma 2.

$$\begin{bmatrix} m-i\\ k\end{bmatrix} q^{i(k+1)} = g(i) - g(i-1) \qquad \textit{where} \qquad g(i) = -\begin{bmatrix} m-i\\ k+1\end{bmatrix} q^{(i+1)(k+1)}. \quad \square$$

Now we can plug into the general formula above and compute

$$D_n = E_n - \sum_{i=1}^{p-1}\sum_{k\ge 0} \begin{bmatrix} n-(p-1)k-i\\ k\end{bmatrix} q^{p\binom{k+1}{2}+i(k+1)} (-1)^k$$
$$= E_n - \sum_{k\ge 0} (-1)^k q^{p\binom{k+1}{2}} \sum_{i=1}^{p-1} \begin{bmatrix} n-(p-1)k-i\\ k\end{bmatrix} q^{i(k+1)}$$
$$= E_n - \sum_{k\ge 0} (-1)^k q^{p\binom{k+1}{2}} \Bigg\{ q^{k+1} \begin{bmatrix} n-(p-1)k\\ k+1\end{bmatrix} - q^{p(k+1)} \begin{bmatrix} n-(p-1)(k+1)\\ k+1\end{bmatrix} \Bigg\}$$
$$= 1 - \sum_{k\ge 0} (-1)^k q^{p\binom{k+1}{2}} q^{k+1} \begin{bmatrix} n-(p-1)k\\ k+1\end{bmatrix},$$

which is the announced formula after a simple change of variable. Note that in the penultimate step the telescoping property of the lemma has been used.

2. Computer algebra proofs

The polynomial families (E_n) and (D_n) give rise to the following study with respect to possible computer proofs. Let us take as input our sum representations of E_n and D_n:

$$\begin{aligned} E_n &= \sum_{k\ge 0} \begin{bmatrix} n-(p-1)k \\ k \end{bmatrix} q^{p\binom{k+1}{2}}(-1)^k, \\ D_n &= \sum_{k\ge 0} \begin{bmatrix} n-(p-1)(k-1) \\ k \end{bmatrix} q^{k+p\binom{k}{2}}(-1)^k. \end{aligned} \tag{1}$$

Then, if p is chosen as a specific positive integer, Riese's package qZeil [6] returns the recurrences $S_n = S_{n-1} - q^n S_{n-p}$ $(n \ge p)$ together with a certificate function Cert for independent verification. Despite the fact that for general "generic" integer parameter p there is no algorithm available, a general pattern can be easily guessed from running the algorithm for $p = 1$, $p = 2$, and $p = 3$, say.

For example, let $F(n,k)$ be the kth summand in our sum representation (1) of E_n, then the recurrence for E_n can be refined to the following statement.

Theorem 1. *For $n \ge p$ and $\delta_k f(n,k) = f(n,k) - f(n,k-1)$, we have*

$$F(n,k) - F(n-1,k) + q^n F(n-p,k) = \delta_k \mathsf{Cert}(n,k) F(n,k), \tag{2}$$

where

$$\mathsf{Cert}(n,k) = q^n \frac{(q^{n-p(k+1)+1})_p}{(q^{n-(p-1)(k+1)})_p}.$$

Proof. After dividing both sides of (2) by $F(n,k)$ the proof reduces to checking equality of rational functions. Namely, note that

$$\begin{aligned} \frac{F(n-1,k)}{F(n,k)} &= \frac{1-q^{n-pk}}{1-q^{n-(p-1)k}}, \\ \frac{F(n,k-1)}{F(n,k)} &= -\frac{q^{pk}}{1-q^k} \frac{(q^{n-pk+1})_p}{(q^{n-(p-1)k+1})_{p-1}}, \end{aligned}$$

and

$$\frac{F(n-p,k)}{F(n,k)} = q^{-n}\mathsf{Cert}(n,k). \qquad \square$$

Analogously, there is a refined version of the recurrence for D_n. The certificate in this case is

$$\mathsf{Cert}(n,k) = q^n \frac{(q^{n-pk})_p}{(q^{n-(p-1)k})_p}.$$

Summarizing, with the sum representation for E_n and D_n in hand, the corresponding recurrences follow immediately by summing both sides of the computer recurrences (2) over all $k \ge 0$.

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