

First Order Predicate Logic

5. Sequent Calculus

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Sequent Calculus: Inference Rules

In addition to the already known sequent calculus rules from propositional logic, we only need to add the rules for quantified formulae:

	Assumptions	Goals
$\forall$	$\frac{\Phi, \forall \gamma, \gamma_{x \rightarrow t} \vdash \Psi}{\Phi, \forall \gamma \vdash \Psi} \quad \text{instantiation}$	$\frac{\Phi \vdash \gamma_{x \rightarrow a}, \Psi}{\Phi \vdash \forall \gamma, \Psi} \vdash \forall \quad \text{arbitrary but fixed}$
$\exists$	$\frac{\Phi, \gamma_{x \rightarrow a} \vdash \Psi}{\Phi, \exists \gamma \vdash \Psi} \quad \text{take such a ...}$	$\frac{\Phi \vdash \gamma_{x \rightarrow t}, \exists \gamma, \Psi}{\Phi \vdash \exists \gamma, \Psi} \vdash \exists \quad \text{witness}$

a : “new” constant (a does not occur in γ, Φ, Ψ)
(Skolem constant)

t : “appropriate” ground term (for any ground term the rule is correct, but only some terms will be appropriate for a successful proof)

Sequent Calculus: Correctness $\forall$

Instantiation

$$\boxed{\frac{\Phi, \forall_x \gamma, \gamma_{x \rightarrow t} \vdash \Psi}{\Phi, \forall_x \gamma \vdash \Psi} \forall \vdash}$$

We know $\forall_x \gamma \models \gamma_{x \rightarrow t}$
for any ground term t .

Prove: $A \Rightarrow B \models ((A \wedge C) \Rightarrow D) \Leftrightarrow ((A \wedge B \wedge C) \Rightarrow D)$

Arbitrary but fixed

$$\boxed{\frac{\Phi \vdash \gamma_{x \rightarrow a}, \Psi}{\Phi \vdash \forall_x \gamma, \Psi} \vdash \forall}$$

$$\begin{aligned} \varphi \Rightarrow \left(\left(\forall_x \gamma \right) \vee \psi \right) &\equiv (\neg \varphi) \vee \left(\left(\forall_a \gamma_{x \rightarrow a} \right) \vee \psi \right) \\ &\equiv \forall_a ((\neg \varphi) \vee (\gamma_{x \rightarrow a} \vee \psi)) \equiv \forall_a (\varphi \Rightarrow (\gamma_{x \rightarrow a} \vee \psi)) \xrightarrow{\text{proof } a} \forall(\mathbb{T}) \equiv \mathbb{T} \end{aligned}$$

Sequent Calculus: Correctness $\exists$

Take such a ...

$$\frac{\Phi, \gamma_{x \rightarrow a} \vdash \Psi}{\Phi, \exists_x \gamma \vdash \Psi} \exists \vdash$$

$$\frac{\frac{\frac{\Phi, \gamma_{x \rightarrow a} \vdash \Psi}{\Phi \vdash \neg \gamma_{x \rightarrow a}, \Psi} \vdash \neg}{\Phi \vdash \forall_x \neg \gamma, \Psi} \vdash \forall}{\Phi, \neg \forall_x \neg \gamma \vdash \Psi} \neg \vdash \equiv \frac{\Phi, \exists_x \gamma \vdash \Psi}{\equiv}$$

Witness

$$\frac{\Phi \vdash \gamma_{x \rightarrow t}, \exists_x \gamma, \Psi}{\Phi \vdash \exists_x \gamma, \Psi} \vdash \exists$$

Exercise 1

Sequent Calculus: Example 1

$$\frac{\frac{\frac{\frac{\overline{Q, P[a], \forall_x P[x] \vdash Q}^a}{P[a] \Rightarrow Q, P[a], \forall_x P[x] \vdash Q} \text{MP}}{P[a] \Rightarrow Q, \forall_x P[x] \vdash Q} \forall \vdash}{\exists_x (P[x] \Rightarrow Q), \forall_x P[x] \vdash Q} \exists \vdash}{\exists_x (P[x] \Rightarrow Q) \vdash \left(\forall_x P[x] \right) \Rightarrow Q} \Rightarrow \vdash$$

Exercise 2: $\forall_x (P[x] \Rightarrow Q) \vdash \left(\exists_x P[x] \right) \Rightarrow Q$

Sequent Calculus: Example 2

$$\begin{array}{c}
 \frac{\overline{P[a] \vdash P[a], Q, \exists_x (P[x] \Rightarrow Q)}^a}{\vdash P[a], P[a] \Rightarrow Q, \exists_x (P[x] \Rightarrow Q)} \vdash \Rightarrow \\
 \frac{\vdash P[a], \exists_x (P[x] \Rightarrow Q)}{\vdash \forall_x P[x], \exists_x (P[x] \Rightarrow Q)} \vdash \exists \\
 \frac{\vdash \forall_x P[x], \exists_x (P[x] \Rightarrow Q)}{\vdash \forall_x P[x] \Rightarrow Q \vdash \exists_x (P[x] \Rightarrow Q)} \vdash \forall \\
 \hline
 \frac{\vdash \forall_x P[x] \Rightarrow Q \vdash \exists_x (P[x] \Rightarrow Q)}{\vdash \forall_x P[x] \Rightarrow Q \vdash \exists_x (P[x] \Rightarrow Q)} \Rightarrow \vdash
 \end{array}$$

$$\begin{array}{c}
 \frac{\overline{Q, P[b] \vdash Q, \exists_x (P[x] \Rightarrow Q)}^a}{Q \vdash P[b] \Rightarrow Q, \exists_x (P[x] \Rightarrow Q)} \vdash \Rightarrow \\
 \frac{Q \vdash P[b] \Rightarrow Q, \exists_x (P[x] \Rightarrow Q)}{Q \vdash \exists_x (P[x] \Rightarrow Q)} \vdash \exists \\
 \hline
 \frac{Q \vdash \exists_x (P[x] \Rightarrow Q)}{Q \vdash \exists_x (P[x] \Rightarrow Q)} \Rightarrow \vdash
 \end{array}$$

Exercise 3: $\left(\exists_x P[x]\right) \Rightarrow Q \vdash \forall_x (P[x] \Rightarrow Q)$

Sequent Calculus: Unit Propagation

Ground atom $\mathcal{A}$ (similar/dual for ground negated atom):

- ▶ match with $\forall$ assumptions, $\exists$ goals, add instances
- ▶ replace $\mathcal{A} \rightarrow \mathbb{T}$ if $\mathcal{A}$ in assumptions, $\mathcal{A} \rightarrow \mathbb{F}$ if $\mathcal{A}$ in goals
- ▶ simplify truth constants

$$\begin{array}{c}
 \frac{}{\mathbb{F}, \forall_x P[x] \vdash} \text{a} \\
 \frac{}{\overline{P[a]}, P[a], \forall_x P[x] \vdash} \text{UP (replace } P[a] \rightarrow \mathbb{F}) \\
 \frac{}{\overline{P[a]}, \forall_x P[x] \vdash} \text{UP (} P[a] \text{ matches } P[x] \text{: instantiate)} \\
 \frac{}{\exists_x \overline{P[x]}, \forall_x P[x] \vdash} \exists \vdash \\
 \frac{}{\exists_x (P[x] \Rightarrow \mathbb{F}), \forall_x P[x] \vdash} \equiv \text{(simplify truth constants)} \\
 \frac{}{\exists_x (P[x] \Rightarrow Q), \forall_x P[x] \vdash} \text{UP (replace } Q \rightarrow \mathbb{F}) \\
 \frac{}{\exists_x (P[x] \Rightarrow Q) \vdash \left(\forall_x P[x] \right) \Rightarrow Q} \Rightarrow \vdash
 \end{array}$$

Exercise 4: use unit propagation: $\left(\forall_x P[x] \right) \Rightarrow Q \vdash \exists_x (P[x] \Rightarrow Q)$

Sequent Calculus: Exercise 1

Witness

$$\boxed{\frac{\Phi \vdash \gamma_{x \rightarrow t}, \exists_x \gamma, \Psi}{\Phi \vdash \exists_x \gamma, \Psi} \vdash \forall}$$

$$\frac{\frac{\frac{\frac{\Phi \vdash \gamma_{x \rightarrow t}, \exists_x \gamma, \Psi}{\Phi \vdash \gamma_{x \rightarrow t}, \neg \forall_x \neg \gamma, \Psi} \equiv}{\Phi, \neg \gamma_{x \rightarrow t}, \forall_x \neg \gamma \vdash \Psi} \vdash \neg}{\Phi, \forall_x \neg \gamma \vdash \Psi} \vdash \forall}{\frac{\Phi \vdash \neg \forall_x \neg \gamma, \Psi}{\Phi \vdash \exists_x \gamma, \Psi} \vdash \neg} \equiv$$

Sequent Calculus: Exercise 2

$$\frac{\frac{\frac{\overline{Q, \forall_x (P[x] \Rightarrow Q), P[a] \vdash Q}^a}{P[a] \Rightarrow Q, \forall_x (P[x] \Rightarrow Q), P[a] \vdash Q} \text{MP}}{\forall_x (P[x] \Rightarrow Q), P[a] \vdash Q} \forall \vdash}{\forall_x (P[x] \Rightarrow Q), \exists_x P[x] \vdash Q} \exists \vdash}{\forall_x (P[x] \Rightarrow Q) \vdash \left(\exists_x P[x] \right) \Rightarrow Q} \Rightarrow \vdash$$

Sequent Calculus: Exercise 3

$$\begin{array}{c}
 \frac{}{P[a] \vdash P[a], \exists_x P[x], Q} \text{a} \\
 \hline
 \frac{P[a] \vdash \exists_x P[x], Q}{\quad} \vdash \exists \quad \frac{}{Q, P[a] \vdash Q} \text{a} \\
 \hline
 \frac{}{(\exists_x P[x]) \Rightarrow Q, P[a] \vdash Q} \Rightarrow \vdash \\
 \hline
 \frac{}{(\exists_x P[x]) \Rightarrow Q \vdash P[a] \Rightarrow Q} \vdash \Rightarrow \\
 \hline
 \frac{}{(\exists_x P[x]) \Rightarrow Q \vdash \forall_x (P[x] \Rightarrow Q)} \vdash \forall
 \end{array}$$

Sequent Calculus: Exercise 4

$$\begin{array}{c}
 \frac{\overline{\vdash \mathbb{T}, \exists_x (P[x] \Rightarrow Q)}^a}{\vdash \mathbb{F} \Rightarrow Q, \exists_x (P[x] \Rightarrow Q)} \equiv \text{(simplify truth constants)} \\
 \frac{\vdash P[a], P[a] \Rightarrow Q, \exists_x (P[x] \Rightarrow Q)}{\vdash P[a], \exists_x (P[x] \Rightarrow Q)} \text{UP (replace } P[a] \rightarrow \mathbb{F} \text{)} \\
 \frac{\vdash P[a], \exists_x (P[x] \Rightarrow Q)}{\vdash \forall_x P[x], \exists_x (P[x] \Rightarrow Q)} \text{UP (} P[a] \text{ matches } P[x]: \text{ instantiate)} \\
 \vdash \forall_x P[x], \exists_x (P[x] \Rightarrow Q) \vdash \forall \\
 \hline
 (\forall_x P[x]) \Rightarrow Q \vdash \exists_x (P[x] \Rightarrow Q)
 \end{array}
 \qquad
 \begin{array}{c}
 \frac{\overline{\vdash \mathbb{T}}^a}{\vdash \exists_x (P[x] \Rightarrow \mathbb{T})} \equiv \\
 \frac{\vdash \exists_x (P[x] \Rightarrow \mathbb{T})}{Q \vdash \exists_x (P[x] \Rightarrow Q)} \text{UP} \\
 \hline
 Q \vdash \exists_x (P[x] \Rightarrow Q) \vdash =
 \end{array}$$

Sequent Calculus: Exercise 4 (continued)

$$\begin{array}{c}
 \frac{\overline{\vdash \mathbb{T}, \exists_x (P[x] \Rightarrow Q)}^a}{\vdash \mathbb{F} \Rightarrow Q, \exists_x (P[x] \Rightarrow Q)} \equiv \\
 \frac{\vdash P[a], P[a] \Rightarrow Q, \exists_x (P[x] \Rightarrow Q)}{\vdash P[a], \exists_x (P[x] \Rightarrow Q)} \text{UP} \\
 \frac{\vdash \forall_x P[x], \exists_x (P[x] \Rightarrow Q)}{\vdash \forall_x P[x], \exists_x (P[x] \Rightarrow Q)} \text{UP} \\
 \frac{\vdash \forall_x P[x], \exists_x (P[x] \Rightarrow Q)}{\vdash \forall_x P[x], \exists_x (P[x] \Rightarrow Q)} \vdash \forall \\
 \frac{\vdash \forall_x P[x], \exists_x (P[x] \Rightarrow Q)}{\vdash \forall_x P[x], \exists_x (P[x] \Rightarrow Q)} \Rightarrow \vdash
 \end{array}$$

$$\begin{array}{c}
 \frac{\overline{\vdash \mathbb{T}}^a}{\vdash \exists_x (P[x] \Rightarrow \mathbb{T})} \equiv (\text{simplify truth constants}) \\
 \frac{Q \vdash \exists_x (P[x] \Rightarrow Q)}{Q \vdash \exists_x (P[x] \Rightarrow Q)} \text{UP (replace } Q \rightarrow \mathbb{T}) \\
 \Rightarrow \vdash
 \end{array}$$

$$\frac{\vdash \forall_x P[x], \exists_x (P[x] \Rightarrow Q)}{Q \vdash \exists_x (P[x] \Rightarrow Q)} \Rightarrow \vdash$$

Unit Propagation: Example 3

$$\begin{array}{c}
 \frac{}{\vdash \mathbb{T}, \exists_x P[x]} \text{ a} \\
 \frac{P[a] \vdash P[a], \exists_x P[x]}{\vdash P[a] \vdash \exists_x P[x]} \text{ UP (replace } P[a] \rightarrow \mathbb{T}) \\
 \frac{}{\vdash P[a] \vdash \exists_x P[x]} \text{ UP (} P[a] \text{ matches } P[x] \text{: instantiate)} \\
 \frac{}{\vdash \neg(\exists_x P[x]), P[a] \vdash} \neg\vdash \\
 \frac{}{\vdash (\exists_x P[x]) \Rightarrow \mathbb{F}, P[a] \vdash} \equiv \text{ (simplify truth constants)} \\
 \frac{}{\vdash (\exists_x P[x]) \Rightarrow Q, P[a] \vdash Q} \text{ UP (replace } Q \rightarrow \mathbb{F}) \\
 \frac{}{\vdash (\exists_x P[x]) \Rightarrow Q \vdash P[a] \Rightarrow Q} \vdash\Rightarrow \\
 \frac{}{\vdash (\exists_x P[x]) \Rightarrow Q \vdash \forall_x (P[x] \Rightarrow Q)} \vdash\forall
 \end{array}$$

Unit Propagation: Example 4

$$\begin{array}{c}
 \frac{}{\mathbb{F}, \forall_x \bar{P}[x] \vdash} \text{a} \\
 \frac{}{\mathbb{T}, \forall_x \bar{P}[x] \vdash} \equiv (\text{simplify truth constants}) \\
 \frac{}{\bar{P}[a], \forall_x \bar{P}[x], P[a] \vdash} \text{UP (replace } P[a] \rightarrow \mathbb{T}) \\
 \frac{}{\forall_x \bar{P}[x], P[a] \vdash} \text{UP (} P[a] \text{ matches } P[x]: \text{ instantiate)} \\
 \frac{}{\forall_x \bar{P}[x], \exists_x P[x] \vdash} \exists \vdash \\
 \frac{}{\forall_x (P[x] \Rightarrow \mathbb{F}), \exists_x P[x] \vdash} \text{Q} \equiv (\text{simplify truth constants}) \\
 \frac{}{\forall_x (P[x] \Rightarrow Q), \exists_x P[x] \vdash} \text{Q} \text{UP (replace } Q \rightarrow \mathbb{F}) \\
 \hline
 \forall_x (P[x] \Rightarrow Q) \vdash (\exists_x P[x]) \Rightarrow Q \vdash \Rightarrow
 \end{array}$$