# Perfect Matchings in Cubic Graphs 

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This talk is dedicated to the memory of


## GERT SABIDUSSI 1929-2022

## Cycles, 2-factors, perfect matchings

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where, for a graph $G, \operatorname{PerMat}(G)$ \# of perfect matchings of $G$.

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Extrem(n) a ladder with a $K_{3,3}$ included at both ends.


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$$
\operatorname{PerMat}(G) \leq \prod_{i=1}^{n} 6^{1 / 3}=6^{n / 3} \text { for cubic bipartite graphs }
$$

A one-to-one correspondence between a binary square matrices and a simple balanced bipartite graphs. Each term of permanent $a_{1 \sigma(1)} \ldots a_{n \sigma(n)}=1$ corresponds to a perfect matching, and vice versa.

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## A formula for PerMat(G)

## Theorem

For a cubic graph $G$ of order $2 n$,

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\operatorname{PerMat}(G)=\mathrm{E}(X), \text { the expected value of } X
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where $X$ is a random variable defined on the set of all 2-colorings $c$ on the edges of $G$, each coloring equally likely, and $X(c)=(-3)^{m_{c}}$, where $m_{c}$ is the number of vertices of $G$ incident in $c$ with three edges of the same color.

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Not feasible for practical calculation but might turn out to have a theoretical value.

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and linear map $\beta$ is given by

$$
\begin{aligned}
\beta: B \otimes B \otimes B & \rightarrow \mathbb{R}, \\
\left|s_{1}\right\rangle \otimes\left|s_{2}\right\rangle \otimes\left|s_{3}\right\rangle & \mapsto \begin{cases}1 & \text { there are exactly two } 0 \text { among } s_{1}, s_{2}, s_{3}, \\
0 & \text { otherwise. }\end{cases}
\end{aligned}
$$

where $s_{1}, s_{2}, s_{3} \in\{0,1\}$, basis of $B \otimes B \otimes B$.

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and the evaluation applies in this case as well.

## Closing remarks

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## Conjecture

Let $G$ be a simple 3-connected cubic graph on $2 n$ vertices, where $n \geq 3$.
Then

$$
\operatorname{PerMat}(G) \leq f_{n}+2 f_{n-1}+2,
$$

with equality attained if and only if $G$ is isomorphic to the circular ladder graph or the Möbius ladder graph

## Verified by a computer for $n \leq 9$.

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Unlike the maximal graphs, the graphs that minimize the number of perfect matchings tend to be complicated

## Erdos - Dongryul Kim



1913-1996


1997 -

