R. Vajda-Z. Kovács (Szeged-Linz): Parametric Root Finding to support discovering geometric inequalities in GeoGebra
(ADG 2021, Hagenberg, September 16)

## Inequality Exploration in planar Euclidean geom.

Create a triangle construction in a DGS and assume that a USER want to explore an INEQUALITY, that is, all the possible ratio between
$m=g_{1} / g_{2}$, where $g_{j}$ can be perimeter, area, circumradius, sum of the medians, etc., for all nondegenerate triangles from a certain class.

We want to support this symbolically! Previous talk:
algebraically as 1st order real quantifier elimination problem (RQE).
We use tarski's (Qepcad) and Mathematica's RQE implementations

## Observations

Because the algebraization in GGBA is coordinate-based, we have several variables in the semialgebraic representations.
When the \#vars is more than 6 , the RQE problem cannot be solved often in a reasonable time ( $\sim 5 \mathrm{sec}$ ) However, the semialgebraic system for certain classes has only FINITELY MANY SOLUTIONS for a fixed $m$, if wlog we fix a triangle side.
For instance, for Isosceles Triangle (IT) or the Right Triangle (RT) classes. $\Rightarrow$
Maybe other methods, approaches can help here to avoid general (full dimensional) RQE and to reduce practical computational time.

## Observations [2]

```
mnf: := {Resolve[Exists[{a, b, c},
            a+b>c^a+c>b^b+c>a^c== b^(a^2+b^2+c^2) == m(ab + b c + ca)], Reals],
    Resolve[Exists[{b}, 2b>1^(1+2b^^2)== m(b^2+2b)],Reals]}
Outf ० = {1\leqm<2, 1\leqm<2}
ln}ff=(\mp@subsup{a}{}{\wedge}2+\mp@subsup{b}{}{\wedge}2+\mp@subsup{c}{}{\wedge}2)/(ab+b c+ca)/.
        {{a->2,b->4,c->4},{a->1,b b 2, c f 2},{a->1/2,b->1,c->1}}
Outf = ={\frac{9}{8},\frac{9}{8},\frac{9}{8}}
mmf f= Solve[(1+2b^^2)== m(b^2 + 2b)/.m -> 9/8]
Outf 0 = {{b }->\frac{4}{7}},{b->2}
```


## Description of the problem

From elementary planar Euclidean geometry:
Consider Inequality Exploration problems from the class of nondegenerate isosceles triangles or the class of right triangles

Out [ : ] =


## Description of the problem [2]

From the algebraic/logical point of view: The EXPLORATION PROBLEM for IT/RT is
not a decision problem, not a SAT/UNSAT problem, but it is very close to that, one free variable $m$ and $n$ existentially bound variables, the nonlinear real algebraic model has Hilbert dimension 1.

We can reduce the RQE problem to finitely many SAT problems, in fact to real root counting (RRC).

## Description of the (new) PRF method

Detect via Gröbner basis computations the "wrong/critical" points of the $m$-parameter space (where the \#sol of the real SAS may change): $O_{\text {crit }} \cup O_{\text {in }} \cup O_{\text {inf }}$ (Computation of MDV via reduction to Elimination)

Decompose the $m$-space into finitely many cells, generate sample for open cells

Solve the Real Root Counting// Real SAT problem and generate a qfree formula based on this. Ref.: [Lazard 2007][Moroz 2006, 2011], [Xia, Hou 2002]

## Description of the (new) PRF method [2]

Existing Implementation : [Maple PRF Package, Maple RegChains]
(not only for 1pm )
Problems/difficulties:
disjunctions, non-strict inequalities, well-behaved systems, orderings.

In fact the IEP for general triangles lead to 2 pms problems, but maybe a recursive classification of the 2D pm space helps!

## A very simple example

For an isosceles triangle, denote the length of the three sides
$\mathrm{AB}, \mathrm{BC}, \mathrm{CA}$, by $\mathrm{a}=1, \mathrm{~b}, \mathrm{c}=\mathrm{b} \quad$ (wlog a=1)

What is the (range of the) ratio of the sum of the squares of the sides $\left(\mathrm{AB}^{2}+\mathrm{BC}^{2}+\mathrm{CA}^{2}=2 b^{2}+1\right)$ and the sum of the products of the sides $\left(\mathrm{ABBC}+\mathrm{ABCA}+\mathrm{BC} \mathrm{CA}=b b+b+b=b^{2}+2 b\right)$ ?

## A very simple example [2]

What is the (range of the) ratio of the sum of the squares of the sides $\left(A B^{2}+B C^{2}+C A^{2}=2 b^{2}+1\right)$ and the sum of the products of the sides $\left(\mathrm{AB} \mathrm{BC}+\mathrm{ABCA}+\mathrm{BC} \mathrm{CA}=b b+b+b=b^{2}+2 b\right)$ ?

As an RQE (NONSAT) problem with one free and one (existentially bound) variable:
Resolve $\left[\exists_{b}\left(2 b-1>0 \wedge 2 b^{2}+1==m\left(b^{2}+2 b\right)\right)\right.$, Reals $]$
Out o ) $1 \leq \mathrm{m}<2$

## A very simple example [3]



## A very simple example [4]

## Reduction via Groebner Basis:

Moroz: 1D Ocrit U Oinequs U Oinfinity $\Longrightarrow$ Induces a complete 1D CAD (open intervals and points)
One typical computation for Ocrit via the partial Jacobian, that is, for detecting the value $m=1$ (first call)
$\ln [\circ]:=$ Flatten [\{Factor [GroebnerBasis [
$\left.\left.\left\{\left(2 b^{\wedge} 2+1\right)-m\left(b^{\wedge} 2+2 b\right), D\left[\left(2 b^{\wedge} 2+1\right)-m\left(b^{\wedge} 2+2 b\right), b\right], m(2 b-1) t+1\right\},\{m\},\{t, b\}\right]\right]$,
Factor [GroebnerBasis [\{(2 b^2+1)-m(b^2+2b), m(2b-1)-u,tu-1\},\{m,u\},\{t,b\}]/. .
$u \rightarrow 0]$, Factor [GroebnerBasis [
$\left.\left.\left.\left.\left\{\left(2 b^{\wedge} 2+u^{\wedge} 2\right)-m\left(b^{\wedge} 2+2 b u\right), t u m(2 b-u)-1, b-1\right\},\{m, u\},\{t, b\}\right] / . u \rightarrow 0\right]\right\}\right]$
Out $:]=\left\{(-1+m)(2+m), m^{2}(-6+5 m),-2+m\right\}$
Out : $]=\left\{\{1\},\left\{\frac{6}{5}\right\},\{2\}\right\}$

## A very simple example [5]

Final Solution $(1 \leq m<2)$ the points and intervals for which the value True assigned
$\ln \left[\circ f=\right.$ AbsoluteTiming [Table[Resolve[Exists[\{b\}, $2 \mathrm{~b}-1>0 \wedge\left(2 \mathrm{~b}^{\wedge} 2+1\right)=m\left(b^{2}+2 b\right)$, Reals], $\{m,\{1 / 2,1,11 / 10,6 / 5,3 / 2,2,3\}\}]]$
out $\cdot J=\{0.013465$, \{False, True, True, True, True, False, False \}\}
$\operatorname{mf} \mid \cdot f=$ Reduce[m==1v1<m<6/5vm==6/5v6/5<m<2,x,Reals]
Out o $\mathrm{I}=1 \leq \mathrm{m}<2$

## A very simple example [6]

## Reduce[

Resolve[Exists[\{v10, v11, v13\}, $(m>0) \wedge(v 11>0) \wedge(v 13>0) \wedge(-4 * v 10 \wedge 2+4 * v 11 \wedge 2-1=0) \wedge$ $\left.\left(-4 * v 10 \wedge 2+4 * v 13^{\wedge} 2-1==0\right) \wedge(-m * v 11 * v 13-m * v 11-m * v 13+v 11 \wedge 2+v 13 \wedge 2+1=0)\right]$, Reals], Reals](*GGBA CB version*)

Out - I= $1 \leq m<2$

## Comparison of Expressions Related to

 Triangle Sides via realgeom, Bottema 1(isosceles triangle, ver. b)

$\cdot a b+b c+c a) \leq\left(a^{2} a+b^{2}+c^{2}\right)<(2) \cdot(a b+b c+c a)$
1

## Motivation

It is a practical work, intuitively it hopes to profit from the reduction of the number of variables (number of CAD-cells) (and from theoretical comp. results).

STAT: For IT/RT, 2-12 (bound) variables.
We have a pre-computed RQE Benchmark sets ( $>100$ test cases)
Can we reach with the PRF method the same or even better results than with RQE?

## Findings

We worked with maple's packages and and our prototype implementation in Mathematica based on a selected BOTTEMA-problem collection (BM 1.1, 1.19, 4.2, 5.1, 5.3, 6.1, 8.1)
All the GGBA generated CB-based IT/RT problems could be treated with PRF but some refinements in our implementation are needed, ongoing work...

```
m| | | = 0.862 ... \leqm< 1/| ToRadicals
Outf. I= - - 
    Inpsas94RBM81pb =
    {{(4 * v16^2-v8^ 2-1), (4 * v17^ 2-4*v8^ 2-1), v18 ^ 2-v8^ 2-1, (v19^ 2-v8^ 2),
        4*v20^2-v8^2-4,(-m*v18-m*v19-m + v16 + v17 + v20)},{m,v16,v17,v18,v19,v20}}
    auxd2d[{Inpsas94,m}]
```

$\left\{\{\mathrm{v} 16, \mathrm{v} 17, \mathrm{v} 18, \mathrm{v} 19, \mathrm{v} 20, \mathrm{v} 8\}, \operatorname{True},\left\{\left(81-72 \mathrm{~m}-232 \mathrm{~m}^{2}-32 \mathrm{~m}^{3}+16 \mathrm{~m}^{4}\right)\right.\right.$
$\left(81+72 m-232 m^{2}+32 m^{3}+16 m^{4}\right)\left(-81+324 m^{2}+36 m^{3}+144 m^{4}+416 m^{5}-64 m^{6}+64 m^{7}\right)$
$\left.\left(81-324 \mathrm{~m}^{2}+36 \mathrm{~m}^{3}-144 \mathrm{~m}^{4}+416 \mathrm{~m}^{5}+64 \mathrm{~m}^{6}+64 \mathrm{~m}^{7}\right)\right\}$,
$\left\{4(-1+m) m^{10}(1+m)(-1+2 m)^{2}(1+2 m)^{2}\left(-3+4 m^{2}\right)^{8}\left(3+4 m^{2}\right)^{8}\right.$
$\left(81+324 m-324 m^{2}-1296 m^{3}+2304 m^{4}-2304 m^{5}+1728 m^{6}+256 m^{8}\right)^{2}$
$\left.\left(81-324 \mathrm{~m}-324 \mathrm{~m}^{2}+1296 \mathrm{~m}^{3}+2304 \mathrm{~m}^{4}+2304 \mathrm{~m}^{5}+1728 \mathrm{~m}^{6}+256 \mathrm{~m}^{8}\right)^{2}\right\}$,
$\left\{\left\{-u^{2}+4 v 16^{2}-v 8^{2}\right\},\left\{-u^{2}+4 v 17^{2}-4 v 8^{2}\right\},\left\{-u^{2}+v 18^{2}-v 8^{2}\right\},\left\{v 19^{2}-v 8^{2}\right\}\right.$,
$\left.\left\{-4 u^{2}+4 v 20^{2}-v 8^{2}\right\},\{-m u+v 16+v 17-m v 18-m v 19+v 20\}\right\}$,
$\{\{m\},\{v 16\},\{v 17\},\{v 18\},\{v 19\},\{v 20\}\},\left\{\left\{-67108864(-1+m) m^{4}(1+m)(-1+2 m)^{2}(1+2 m)^{2}\right\}\right.$,
$\left\{-1048576(-1+m) m^{4}(1+m)(-1+2 m)^{2}(1+2 m)^{2}\right\},\left\{-16384(-1+m) m^{4}(1+m)(-1+2 m)^{2}(1+2 m)^{2}\right\}$,
$\left\{16384(-1+m) m^{4}(1+m)(-1+2 m)^{2}(1+2 m)^{2}\right\},\left\{-1048576(-1+m) m^{4}(1+m)(-1+2 m)^{2}(1+2 m)^{2}\right\}$,

$$
\begin{aligned}
& \left\{\left\{\left\{m \rightarrow \frac{27}{226}\right\}\right\},\left\{\left\{m \rightarrow \frac{96153}{214879}\right\}\right\},\left\{\left\{m \rightarrow \frac{893}{1936}\right\}\right\},\left\{\left\{m \rightarrow \frac{1092}{2131}\right\}\right\},\left\{\left\{m \rightarrow \frac{35}{57}\right\}\right\},\right. \\
& \left.\left\{\left\{m \rightarrow \frac{452}{553}\right\}\right\},\left\{\left\{m \rightarrow \frac{21851}{25317}\right\}\right\},\left\{\left\{m \rightarrow \frac{679}{754}\right\}\right\},\left\{\left\{m \rightarrow \frac{10}{7}\right\}\right\},\left\{\left\{m \rightarrow \frac{68}{21}\right\}\right\},\{\{m \rightarrow 32\}\}\right\}, \\
& \text { \{False, False, False, False, False, False, True, True, False, False, False\}, } \\
& \left\{\text { False, False, False, False, False, False, } \sqrt{ } 0.862 \ldots<m<\frac{\sqrt{3}}{2}\right. \text {, } \\
& \left.\frac{\sqrt{3}}{2}<m<1 \text {, False, False, False }\right\},\{\text { False, False, False, False, False, }
\end{aligned}
$$

## Conclusion

If GB and (nonlinear) real SAT or RRC are implemented and they are fast, it COULD be a viable approach instead of the general RQE.
Educational applications all the sub-algorithms should be implemented in a free software (GB $\longrightarrow$ Giac, SAT $\longrightarrow$ tarski, SMT-RAT,..., WS?)

The exploration problems for a GENERAL triangles are not 1D problems. MDV in a 2 D space: 2D generic CAD, also recursive analysis of curves?
Discussion: Any suggestion? SEE GT for $m=\left(A B^{2}+B C^{2}+C A^{2}\right) /(A B+B C+C A)$


## References

[Moroz 2011] Properness defects of projection and minimal discriminant variety, Journal of Symbolic Computation 46(10), 1139-1157, 2011.
[Liang-Gerhard-Jeffrey-Moroz 2009] A package for solving parametric polynomial systems, ACM Communications in Computer Algebra 169(43), 2009.
[realgeom] GeoGebra and the realgeom Reasoning Tool, CEUR Workshop Proceedings Vol. 2752, PAAR+SC-Square Workshop, Paris, France, 204-219, 2020.

