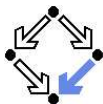
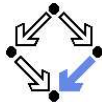


Verifying Concurrent Systems with PVS

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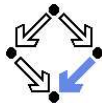
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<http://www.risc.uni-linz.ac.at>



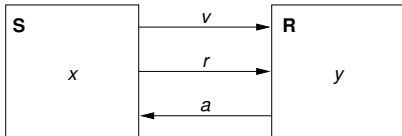


1. A Bit Transmission Protocol

2. A Client/Server System



A Bit Transmission Protocol



```
var x, y  
var v := 0, r := 0, a := 0
```

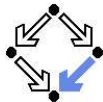
S: loop

```
  choose x ∈ {0, 1}    ||  
  1 : v, r := x, 1  
  2 : wait a = 1  
    r := 0  
  3 : wait a = 0
```

R: loop

```
  1 : wait r = 1  
    y, a := v, 1  
  2 : wait r = 0  
    a := 0
```

Transmit a bit through a wire.



A (Simplified) Model of the Protocol

$$\text{State} := PC^2 \times (\mathbb{N}_2)^5$$

$$I(p, q, x, y, v, r, a) :\Leftrightarrow p = q = 1 \wedge x \in \mathbb{N}_2 \wedge v = r = a = 0.$$

$$R(\langle p, q, x, y, v, r, a \rangle, \langle p', q', x', y', v', r', a' \rangle) :\Leftrightarrow \\ S1(\dots) \vee S2(\dots) \vee S3(\dots) \vee R1(\dots) \vee R2(\dots).$$

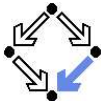
$$S1(\langle p, q, x, y, v, r, a \rangle, \langle p', q', x', y', v', r', a' \rangle) :\Leftrightarrow \\ p = 1 \wedge p' = 2 \wedge v' = x \wedge r' = 1 \wedge \\ q' = q \wedge x' = x \wedge y' = y \wedge v' = v \wedge a' = a.$$

$$S2(\langle p, q, x, y, v, r, a \rangle, \langle p', q', x', y', v', r', a' \rangle) :\Leftrightarrow \\ p = 2 \wedge p' = 3 \wedge a = 1 \wedge r' = 0 \wedge \\ q' = q \wedge x' = x \wedge y' = y \wedge v' = v \wedge a' = a.$$

$$R1(\langle p, q, x, y, v, r, a \rangle, \langle p', q', x', y', v', r', a' \rangle) :\Leftrightarrow \\ p = 3 \wedge p' = 1 \wedge a = 0 \wedge x' \in \mathbb{N}_2 \wedge \\ q' = q \wedge y' = y \wedge v' = v \wedge r' = r \wedge a' = a.$$

$$R2(\langle p, q, x, y, v, r, a \rangle, \langle p', q', x', y', v', r', a' \rangle) :\Leftrightarrow \\ q = 1 \wedge q' = 2 \wedge r = 1 \wedge y' = v \wedge a' = 1 \wedge \\ p' = p \wedge x' = x \wedge v' = v \wedge r' = r.$$

$$R3(\langle p, q, x, y, v, r, a \rangle, \langle p', q', x', y', v', r', a' \rangle) :\Leftrightarrow \\ q = 2 \wedge q' = 1 \wedge r = 0 \wedge a' = 0 \wedge \\ p' = p \wedge x' = x \wedge y' = y \wedge v' = v \wedge r' = r.$$



A Verification Task

$$\langle I, R \rangle \models \Box(q = 1 \Rightarrow y = x)$$

$$\text{Invariant}(p, \dots) \Rightarrow (q = 1 \Rightarrow y = x)$$

$$I(p, \dots) \Rightarrow \text{Invariant}(p, \dots)$$

$$R(\langle p, \dots \rangle, \langle p', \dots \rangle) \wedge \text{Invariant}(p, \dots) \Rightarrow \text{Invariant}(p', \dots)$$

$$\text{Invariant}(p, q, x, y, v, r, a) :\Leftrightarrow$$

$$(p = 1 \vee p = 2 \vee p = 3) \wedge (q = 1 \vee q = 2) \wedge$$

$$(x = 0 \vee x = 1) \wedge (v = 0 \vee v = 1) \wedge (r = 0 \vee r = 1) \wedge (a = 0 \vee a = 1) \wedge$$

$$(p = 1 \Rightarrow q = 1 \wedge r = 0 \wedge a = 0) \wedge$$

$$(p = 2 \Rightarrow r = 1) \wedge$$

$$(p = 3 \Rightarrow r = 0) \wedge$$

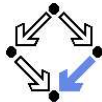
$$(q = 1 \Rightarrow a = 0) \wedge$$

$$(q = 2 \Rightarrow (p = 2 \vee p = 3) \wedge a = 1 \wedge y = x) \wedge$$

$$(r = 1 \Rightarrow p = 2 \wedge v = x)$$

The invariant captures the essence of the protocol.

The Verification Task in PVS



```
protocol: THEORY
BEGIN
```

```
  p, q, x, y, v, r, a: nat
  p0, q0, x0, y0, v0, r0, a0: nat
```

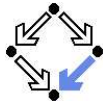
```
S1: bool =
  p = 1 AND p0 = 2 AND v0 = x AND r0 = 1 AND
  q0 = q AND x0 = x AND y0 = y AND v0 = v AND a0 = a
```

```
S2: bool =
  p = 2 AND p0 = 3 AND a = 1 AND r0 = 0 AND
  q0 = q AND x0 = x AND y0 = y AND v0 = v AND a0 = a
```

```
S3: bool =
  p = 3 AND p0 = 1 AND a = 0 AND (x0 = 0 OR x0 = 1) AND
  q0 = q AND y0 = y AND v0 = v AND r0 = r AND a0 = a
```

```
R1: bool =
  q = 1 AND q0 = 2 AND r = 1 AND y0 = v AND a0 = 1 AND
  p0 = p AND x0 = x AND v0 = v AND r0 = r
```

```
R2: bool =
  q = 2 AND q0 = 1 AND r = 0 AND a0 = 0 AND
  p0 = p AND x0 = x AND y0 = y AND v0 = v AND r0 = r
```



The Verification Task in PVS (Contd)

Init: bool =

$p = 1 \text{ AND } q = 1 \text{ AND } (x = 0 \text{ OR } x = 1) \text{ AND}$

$v = 0 \text{ AND } r = 0 \text{ AND } a = 0$

Step: bool =

$S1 \text{ OR } S2 \text{ OR } S3 \text{ OR } R1 \text{ OR } R2$

Property: bool =

$q = 2 \Rightarrow y = x$

Invariant(p, q, x, y, v, r, a: nat): bool =

$(p = 1 \text{ OR } p = 2 \text{ OR } p = 3) \text{ AND}$

$(q = 1 \text{ OR } q = 2) \text{ AND}$

$(x = 0 \text{ OR } x = 1) \text{ AND}$

$(v = 0 \text{ OR } v = 1) \text{ AND}$

$(r = 0 \text{ OR } r = 1) \text{ AND}$

$(a = 0 \text{ OR } a = 1) \text{ AND}$

$(p = 1 \Rightarrow q = 1 \text{ AND } r = 0 \text{ AND } a = 0) \text{ AND}$

$(p = 2 \Rightarrow r = 1) \text{ AND}$

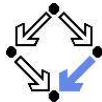
$(p = 3 \Rightarrow r = 0) \text{ AND}$

$(q = 1 \Rightarrow a = 0) \text{ AND}$

$(q = 2 \Rightarrow (p = 2 \text{ OR } p = 3) \text{ AND } a = 1 \text{ AND } y = x) \text{ AND}$

$(r = 1 \Rightarrow (p = 2 \text{ AND } v = x))$

The Verification Task in PVS (Contd'2)



VC0: THEOREM

Invariant(p, q, x, y, v, r, a) => Property

VC1: THEOREM

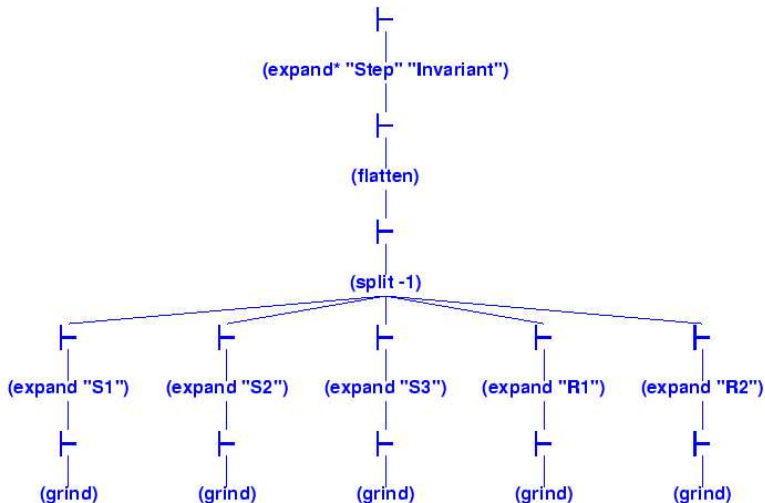
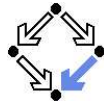
Init => Invariant(p, q, x, y, v, r, a)

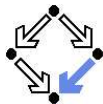
VC2: THEOREM

Step AND Invariant(p, q, x, y, v, r, a) =>
Invariant(p0, q0, x0, y0, v0, r0, a0)

END protocol

The Proof in PVS

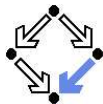




1. A Bit Transmission Protocol

2. A Client/Server System

The Client



Client system $C_i = \langle IC_i, RC_i \rangle$.

State := $PC \times \mathbb{N}_2 \times \mathbb{N}_2$.

Int := $\{R_i, S_i, C_i\}$.

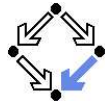
$IC_i(pc, request, answer) :\Leftrightarrow$
 $pc = R \wedge request = 0 \wedge answer = 0$.

$RC_i(I, \langle pc, request, answer \rangle,$
 $\langle pc', request', answer' \rangle) :\Leftrightarrow$
 $(I = R_i \wedge pc = R \wedge request = 0 \wedge$
 $pc' = S \wedge request' = 1 \wedge answer' = answer) \vee$
 $(I = S_i \wedge pc = S \wedge answer \neq 0 \wedge$
 $pc' = C \wedge request' = request \wedge answer' = 0) \vee$
 $(I = C_i \wedge pc = C \wedge request = 0 \wedge$
 $pc' = R \wedge request' = 1 \wedge answer' = answer) \vee$

$(I = \overline{REQ}_i \wedge request \neq 0 \wedge$
 $pc' = pc \wedge request' = 0 \wedge answer' = answer) \vee$
 $(I = \overline{ANS}_i \wedge$
 $pc' = pc \wedge request' = request \wedge answer' = 1).$

```
Client(ident):  
  param ident  
  begin  
    loop  
      ...  
    R: sendRequest()  
    S: receiveAnswer()  
    C: // critical region  
      ...  
      sendRequest()  
    endloop  
  end Client
```

The Server



Server system $S = \langle IS, RS \rangle$.

$State := (\mathbb{N}_3)^3 \times (\{1, 2\} \rightarrow \mathbb{N}_2)^2$.

$Int := \{D1, D2, F, A1, A2, W\}$.

$IS(given, waiting, sender, rbuffer, sbuffer) :\Leftrightarrow$
 $given = waiting = sender = 0 \wedge$
 $rbuffer(1) = rbuffer(2) = sbuffer(1) = sbuffer(2) = 0$.

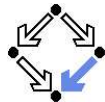
$RS(I, \langle given, waiting, sender, rbuffer, sbuffer \rangle,$
 $\langle given', waiting', sender', rbuffer', sbuffer' \rangle) :\Leftrightarrow$
 $\exists i \in \{1, 2\} :$
 $(I = D_i \wedge sender = 0 \wedge rbuffer(i) \neq 0 \wedge$
 $sender' = i \wedge rbuffer'(i) = 0 \wedge$
 $U(given, waiting, sbuffer) \wedge$
 $\forall j \in \{1, 2\} \setminus \{i\} : U_j(rbuffer)) \vee$
...

$U(x_1, \dots, x_n) :\Leftrightarrow x'_1 = x_1 \wedge \dots \wedge x'_n = x_n$.

$U_j(x_1, \dots, x_n) :\Leftrightarrow x'_1(j) = x_1(j) \wedge \dots \wedge x'_n(j) = x_n(j)$.

```
Server:
  local given, waiting, sender
  begin
    given := 0; waiting := 0
  loop
    D: sender := receiveRequest()
      if sender = given then
        if waiting = 0 then
          F: given := 0
        else
          A1: given := waiting;
              waiting := 0
              sendAnswer(given)
        endif
      elsif given = 0 then
        A2: given := sender
            sendAnswer(given)
      else
        W: waiting := sender
      endif
    endloop
  end Server
```

The Server (Contd)



...

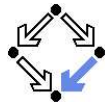
$$(I = F \wedge sender \neq 0 \wedge sender = given \wedge waiting = 0 \wedge given' = 0 \wedge sender' = 0 \wedge U(waiting, rbuffer, sbuffer)) \vee$$
$$(I = A1 \wedge sender \neq 0 \wedge sbuffer(waiting) = 0 \wedge sender = given \wedge waiting \neq 0 \wedge given' = waiting \wedge waiting' = 0 \wedge sbuffer'(waiting) = 1 \wedge sender' = 0 \wedge U(rbuffer) \wedge \forall j \in \{1, 2\} \setminus \{waiting\} : U_j(sbuffer)) \vee$$
$$(I = A2 \wedge sender \neq 0 \wedge sbuffer(sender) = 0 \wedge sender \neq given \wedge given = 0 \wedge given' = sender \wedge sbuffer'(sender) = 1 \wedge sender' = 0 \wedge U(waiting, rbuffer) \wedge \forall j \in \{1, 2\} \setminus \{sender\} : U_j(sbuffer)) \vee$$

...

Server:

```
local given, waiting, sender
begin
  given := 0; waiting := 0
  loop
D: sender := receiveRequest()
    if sender = given then
      if waiting = 0 then
F:       given := 0
        else
A1:      given := waiting;
          waiting := 0
          sendAnswer(given)
        endif
      elsif given = 0 then
A2:      given := sender
          sendAnswer(given)
        else
W:       waiting := sender
        endif
      endif
    endloop
end Server
```

The Server (Contd'2)



...

$$(I = W \wedge sender \neq 0 \wedge sender \neq given \wedge given \neq 0 \wedge$$
$$waiting' := sender \wedge sender' = 0 \wedge$$
$$U(given, rbuffer, sbuffer)) \vee$$

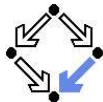
$\exists i \in \{1, 2\} :$

$$(I = REQ_i \wedge rbuffer'(i) = 1 \wedge$$
$$U(given, waiting, sender, sbuffer) \wedge$$
$$\forall j \in \{1, 2\} \setminus \{i\} : U_j(rbuffer)) \vee$$
$$(I = \overline{ANS}_i \wedge sbuffer(i) \neq 0 \wedge$$
$$sbuffer'(i) = 0 \wedge$$
$$U(given, waiting, sender, rbuffer) \wedge$$
$$\forall j \in \{1, 2\} \setminus \{i\} : U_j(sbuffer)).$$

Server:

```
local given, waiting, sender
begin
  given := 0; waiting := 0
  loop
D: sender := receiveRequest()
    if sender = given then
      if waiting = 0 then
F:       given := 0
        else
A1:      given := waiting;
          waiting := 0
          sendAnswer(given)
        endif
      elsif given = 0 then
A2:      given := sender
          sendAnswer(given)
        else
W:       waiting := sender
        endif
      endloop
end Server
```

The Composed System

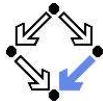


$$\text{State} := (\{1, 2\} \rightarrow PC) \times (\{1, 2\} \rightarrow \mathbb{N}_2)^2 \times (\mathbb{N}_3)^2 \times (\{1, 2\} \rightarrow \mathbb{N}_2)^2$$

$$I(\langle pc, request, answer, given, waiting, sender, rbuffer, sbuffer \rangle) :\Leftrightarrow \\ \forall i \in \{1, 2\} : IC(pc_i, request_i, answer_i) \wedge \\ IS(given, waiting, sender, rbuffer, sbuffer)$$

$$R(\langle pc, request, answer, given, waiting, sender, rbuffer, sbuffer \rangle, \\ \langle pc', request', answer', given', waiting', sender', rbuffer', sbuffer' \rangle) :\Leftrightarrow \\ (\exists i \in \{1, 2\} : RC_{local}(\langle pc_i, request_i, answer_i \rangle, \langle pc'_i, request'_i, answer'_i \rangle) \wedge \\ \langle given, waiting, sender, rbuffer, sbuffer \rangle = \\ \langle given', waiting', sender', rbuffer', sbuffer' \rangle) \vee \\ (RS_{local}(\langle given, waiting, sender, rbuffer, sbuffer \rangle, \\ \langle given', waiting', sender', rbuffer', sbuffer' \rangle) \wedge \\ \forall i \in \{1, 2\} : \langle pc_i, request_i, answer_i \rangle = \langle pc'_i, request'_i, answer'_i \rangle) \vee \\ (\exists i \in \{1, 2\} : External(i, \langle request_i, answer_i, rbuffer, sbuffer \rangle, \\ \langle request'_i, answer'_i, rbuffer', sbuffer' \rangle) \wedge \\ pc = pc' \wedge \langle sender, waiting, given \rangle = \langle sender', waiting', given' \rangle)$$

The Verification Task



$$\langle I, R \rangle \models \Box \neg (pc_1 = C \wedge pc_2 = C)$$

Invariant(*pc*, *request*, *answer*, *sender*, *given*, *waiting*, *rbuffer*, *sbuffer*) : \Leftrightarrow

$\forall i \in \{1, 2\} :$

$$(pc(i) = C \vee sbuffer(i) = 1 \vee answer(i) = 1 \Rightarrow$$

$$given = i \wedge$$

$$\forall j : j \neq i \Rightarrow pc(j) \neq C \wedge sbuffer(j) = 0 \wedge answer(j) = 0) \wedge$$

$$(pc(i) = R \Rightarrow$$

$$sbuffer(i) = 0 \wedge answer(i) = 0 \wedge$$

$$(i = given \Leftrightarrow request(i) = 1 \vee rbuffer(i) = 1 \vee sender = i) \wedge$$

$$(request(i) = 0 \vee rbuffer(i) = 0)) \wedge$$

$$(pc(i) = S \Rightarrow$$

$$(sbuffer(i) = 1 \vee answer(i) = 1 \Rightarrow$$

$$request(i) = 0 \wedge rbuffer(i) = 0 \wedge sender \neq i) \wedge$$

$$(i \neq given \Rightarrow$$

$$request(i) = 0 \vee rbuffer(i) = 0)) \wedge$$

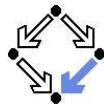
$$(pc(i) = C \Rightarrow$$

$$request(i) = 0 \wedge rbuffer(i) = 0 \wedge sender \neq i \wedge$$

$$sbuffer(i) = 0 \wedge answer(i) = 0) \wedge$$

...

The Verification Task (Contd)

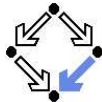


...

$$(sender = 0 \wedge (request(i) = 1 \vee rbuffer(i) = 1) \Rightarrow$$
$$sbuffer(i) = 0 \wedge answer(i) = 0) \wedge$$
$$(sender = i \Rightarrow$$
$$(waiting \neq i) \wedge$$
$$(sender = given \wedge pc(i) = R \Rightarrow$$
$$request(i) = 0 \wedge rbuffer(i) = 0) \wedge$$
$$(pc(i) = S \wedge i \neq given \Rightarrow$$
$$request(i) = 0 \wedge rbuffer(i) = 0) \wedge$$
$$(pc(i) = S \wedge i = given \Rightarrow$$
$$request(i) = 0 \vee rbuffer(i) = 0)) \wedge$$
$$(waiting = i \Rightarrow$$
$$given \neq i \wedge pc_i = S \wedge request_i = 0 \wedge rbuffer(i) = 0 \wedge$$
$$sbuffer_i = 0 \wedge answer(i) = 0) \wedge$$
$$(sbuffer(i) = 1 \Rightarrow$$
$$answer(i) = 0 \wedge request(i) = 0 \wedge rbuffer(i) = 0)$$

As usual, the invariant has been elaborated in the course of the proof.

The Verification Task in PVS



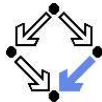
```
clientServer: THEORY
BEGIN

% client indices and program counter constants
Index : TYPE+ = { x: nat | x = 1 OR x = 2 } CONTAINING 1
Index0: TYPE+ = { x: nat | x < 3 } CONTAINING 0
PC: TYPE+ = { R, S, C }

% client states
pc, pc0: [ Index -> PC ]
request, request0: [ Index -> bool ]
answer, answer0: [ Index -> bool ]

% server states
given, given0: Index0
waiting, waiting0: Index0
sender, sender0: Index0
rbuffer, rbuffer0: [ Index -> bool ]
sbuffer, sbuffer0: [ Index -> bool ]
```

The Verification Task in PVS (Contd)



```
i, j: VAR Index
```

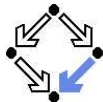
```
% -----  
% initial state condition  
% -----
```

```
IC(pc: PC, request: bool, answer: bool): bool =  
  pc = R AND request = FALSE AND answer = FALSE
```

```
IS(given: Index0, waiting: Index0, sender: Index0,  
  rbuffer: [ Index -> bool ], sbuffer: [ Index -> bool ]): bool =  
  given = 0 AND waiting = 0 AND sender = 0 AND  
  (FORALL i: rbuffer(i) = FALSE AND sbuffer(i) = FALSE)
```

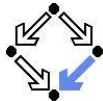
```
Initial: bool =  
  (FORALL i: IC(pc(i), request(i), answer(i))) AND  
  IS(given, waiting, sender, rbuffer, sbuffer)
```

The Verification Task in PVS (Contd'2)



```
% -----  
% transition relation  
% -----  
RC(pc: PC, request: bool, answer: bool,  
   pc0: PC, request0: bool, answer0: bool): bool =  
  (pc = R AND request = FALSE AND  
   pc0 = S AND request0 = TRUE AND answer0 = answer) OR  
  (pc = S AND answer = TRUE AND  
   pc0 = C AND request0 = request AND answer0 = FALSE) OR  
  (pc = C AND request = FALSE AND  
   pc0 = R and request0 = TRUE AND answer0 = answer)  
  
RS(given: Index0, waiting: Index0, sender: Index0,  
   rbuffer: [ Index -> bool ], sbuffer: [ Index -> bool ],  
   given0: Index0, waiting0: Index0, sender0: Index0,  
   rbuffer0: [ Index -> bool ], sbuffer0: [ Index -> bool ]): bool =  
(EXISTS i:  
  sender = 0 AND rbuffer(i) = TRUE AND  
  sender0 = i AND rbuffer0(i) = FALSE AND  
  given = given0 AND waiting = waiting0 AND sbuffer = sbuffer0 AND  
  FORALL j: j /= i => rbuffer(j) = rbuffer0(j)) OR
```

The Verification Task in PVS (Contd'3)



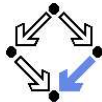
```
(sender /= 0 AND sender = given AND waiting = 0 AND
  given0 = 0 AND sender0 = 0 AND
  waiting = waiting0 AND rbuffer = rbuffer0 AND sbuffer = sbuffer0) OR
```

```
(sender /= 0 AND
  sender = given AND waiting /= 0 AND
  sbuffer(waiting) = FALSE AND % change order for type-checking
  given0 = waiting AND waiting0 = 0 AND
  sbuffer0(waiting) = TRUE AND sender0 = 0 AND
  rbuffer = rbuffer0 AND
  (FORALL j: j /= waiting => sbuffer(j) = sbuffer0(j))) OR
```

```
(sender /= 0 AND sbuffer(sender) = FALSE AND
  sender /= given AND given = 0 AND
  given0 = sender AND
  sbuffer0(sender) = TRUE AND sender0 = 0 AND
  waiting = waiting0 AND rbuffer = rbuffer0 AND
  (FORALL j: j /= sender => sbuffer(j) = sbuffer0(j))) OR
```

```
(sender /= 0 AND sender /= given AND given /= 0 AND
  waiting0 = sender AND sender0 = 0 AND
  given = given0 AND rbuffer = rbuffer0 AND sbuffer = sbuffer0)
```

The Verification Task in PVS (Contd'4)

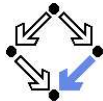


```
External(i: Index,  
  pc: PC, request: bool, answer: bool,  
  pc0: PC, request0: bool, answer0: bool,  
  given: Index0, waiting: Index0, sender: Index0,  
    rbuffer: [ Index -> bool ], sbuffer: [ Index -> bool ],  
  given0: Index0, waiting0: Index0, sender0: Index0,  
    rbuffer0: [ Index -> bool ], sbuffer0: [ Index -> bool ]): bool =
```

```
(request = TRUE AND  
  pc0 = pc AND request0 = FALSE AND answer0 = answer AND  
  rbuffer0(i) = TRUE AND  
  given = given0 AND waiting = waiting0 AND sender = sender0 AND  
  sbuffer = sbuffer0 AND  
  (FORALL j: j /= i => rbuffer(j) = rbuffer0(j))) OR
```

```
(pc0 = pc AND request0 = request AND answer0 = TRUE AND  
  sbuffer(i) = TRUE AND sbuffer0(i) = FALSE AND  
  given = given0 AND waiting = waiting0 AND sender = sender0 AND  
  rbuffer = rbuffer0 AND  
  (FORALL j: j /= i => sbuffer(j) = sbuffer0(j)))
```

The Verification Task in PVS (Contd'5)

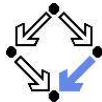


```
Next: bool =
  ((EXISTS i: RC(pc (i), request (i), answer (i),
                pc0(i), request0(i), answer0(i)) AND
   (FORALL j: j /= i =>
    pc(j) = pc0(j) AND request(j) = request0(j) AND
    answer(j) = answer0(j))) AND
   given = given0 AND waiting = waiting0 AND sender = sender0 AND
   rbuffer = rbuffer0 AND sbuffer = sbuffer0) OR

  (RS(given, waiting, sender, rbuffer, sbuffer,
      given0, waiting0, sender0, rbuffer0, sbuffer0) AND
   (FORALL j: pc(j) = pc0(j) AND request(j) = request0(j) AND
    answer(j) = answer0(j))) OR

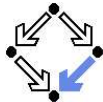
  (EXISTS i:
   External(i, pc (i), request (i), answer (i),
            pc0(i), request0(i), answer0(i),
            given, waiting, sender, rbuffer, sbuffer,
            given0, waiting0, sender0, rbuffer0, sbuffer0) AND
   (FORALL j: j /= i =>
    pc(j) = pc0(j) AND request(j) = request0(j) AND
    answer(j) = answer0(j)))
```

The Verification Task in PVS (Contd'6)



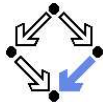
```
% -----  
% invariant  
% -----  
Invariant(pc: [Index->PC], request: [Index -> bool],  
          answer: [Index -> bool],  
          given: Index0, waiting: Index0, sender: Index0,  
          rbuffer: [Index -> bool], sbuffer: [Index->bool]): bool =  
FORALL i:  
  (pc(i) = C OR sbuffer(i) = TRUE OR answer(i) = TRUE =>  
   given = i AND  
   FORALL j: j /= i =>  
     pc(j) /= C AND  
     sbuffer(j) = FALSE AND answer(j) = FALSE) AND  
(pc(i) = R =>  
  sbuffer(i) = FALSE AND answer(i) = FALSE AND  
  (i /= given =>  
   request(i) = FALSE AND rbuffer(i) = FALSE AND sender /= i)  
  (i = given =>  
   request(i) = TRUE OR rbuffer(i) = TRUE OR sender = i) AND  
  (request(i) = FALSE OR rbuffer(i) = FALSE)) AND
```


The Verification Task in PVS (Contd'7)



```
(pc(i) = S =>
  (sbuffer(i) = TRUE OR answer(i) = TRUE =>
    request(i) = FALSE AND rbuffer(i) = FALSE AND sender /= i) AND
  (i /= given =>
    request(i) = FALSE OR rbuffer(i) = FALSE)) AND
(pc(i) = C =>
  request(i) = FALSE AND rbuffer(i) = FALSE AND sender /= i AND
  sbuffer(i) = FALSE AND answer(i) = FALSE) AND
(sender = 0 AND (request(i) = TRUE OR rbuffer(i) = TRUE) =>
  sbuffer(i) = FALSE AND answer(i) = FALSE) AND
(sender = i =>
  (sender = given AND pc(i) = R =>
    request(i) = FALSE and rbuffer(i) = FALSE) AND
  (waiting /= i) AND
  (pc(i) = S AND i /= given =>
    request(i) = FALSE AND rbuffer(i) = FALSE) AND
  (pc(i) = S AND i = given =>
    request(i) = FALSE OR rbuffer(i) = FALSE)) AND
```

The Verification Task in PVS (Contd'8)



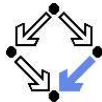
```
(waiting = i =>
  given /= i AND
  pc(waiting) = S AND
  request(waiting) = FALSE AND rbuffer(waiting) = FALSE AND
  sbuffer(waiting) = FALSE AND answer(waiting) = FALSE) AND
(sbuffer(i) = TRUE =>
  answer(i) = FALSE AND request(i) = FALSE AND rbuffer(i) = FALSE)
```

```
% -----
% mutual exclusion proof
% -----
```

Mutex: THEOREM

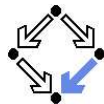
```
Invariant(pc, request, answer,
  given, waiting, sender, rbuffer, sbuffer) =>
NOT (pc(1) = C AND pc(2) = C)
```

The Verification Task in PVS (Contd'9)



```
% -----  
% invariance proof  
% -----  
Inv1: THEOREM  
  Initial =>  
  Invariant(pc, request, answer,  
    given, waiting, sender, rbuffer, sbuffer)  
  
Inv2: THEOREM  
  Invariant(pc, request, answer,  
    given, waiting, sender, rbuffer, sbuffer) AND Next =>  
  Invariant(pc0, request0, answer0,  
    given0, waiting0, sender0, rbuffer0, sbuffer0)  
  
END clientServer
```

The Proof in PVS

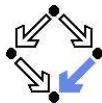


Proofs that the system invariant implies the mutual exclusion property and that the initial condition implies the invariant.

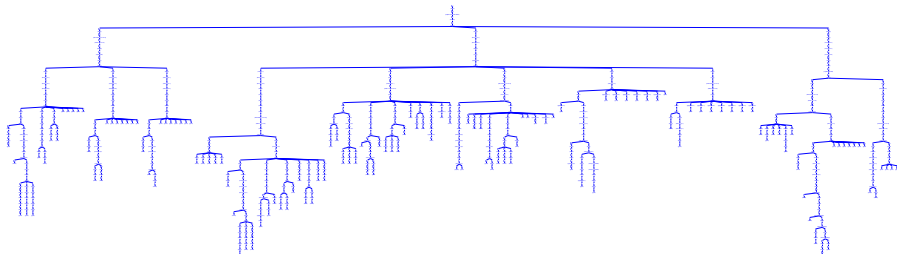
```
┆
┆ (expand* "Invariant0")
┆
┆ (flatten)
┆
┆ (assert)
┆
┆ (inst-cp -1 "1")
┆
┆ (inst-cp -1 "2")
┆
┆ (assert)
┆
┆ (grind)
```

```
┆
┆ (expand* "Initial" "Invariant0")
┆
┆ (expand* "IS")
┆
┆ (expand* "IC")
┆
┆ (grind)
```

The Proof in PVS



Proof that every system transition preserves the invariant.



- 10 subproofs, one for each transition.
 - Three from client, five from server, two from communication system.
 - Download and investigate from course Web site.

Only with computer support, verification proofs become manageable.